# Game Playing Part 2 Alpha-Beta Pruning

# Yingyu Liang yliang@cs.wisc.edu Computer Sciences Department University of Wisconsin, Madison

[based on slides from A. Moore, C. Dyer, J. Skrentny, Jerry Zhu]

#### **Review: Game Theoretic Value**



The bottom-up approach: needs the whole game tree. Too much space! The minimax algorithm: a variant of DFS to save space.

# **Review: Minimax algorithm**

```
function Max-Value(s) inputs:
```

s: current state in game, Max about to play output: *best-score (for Max) available from s* 

```
if ( s is a terminal state )
then return ( terminal value of s )
else
```

```
\alpha := -\infty
for each s' in Succ(s)
\alpha := \max(\alpha, Min-value(s'))
```

return α

```
function Min-Value(s)
output: best-score (for Min) available from s
if ( s is a terminal state )
```

```
then return (terminal value of s) else
```

```
\begin{array}{l} \beta := \infty \\ \text{for each s' in Succs(s)} \\ \beta := \min(\beta, \text{Max-value(s')}) \\ \text{return } \beta \end{array}
```

Time complexity? O(b<sup>m</sup>) ← bad

Space complexity?
 O(bm)







The execution on the terminal nodes is omitted.





















#### alpha-beta pruning

# Gives the same game theoretic values as minimax, but prunes part of the game tree.

"If you have an idea that is surely bad, don't take the time to see how truly awful it is." -- Pat Winston



slide 18



The subtree below G is omitted



is omitted



#### **Alpha-Beta Motivation Summary**



- After returning from A, Max can get at least 100 at S
- After returning from F, Max can get at most 20 at B
- At this point, Max losts interest in B
- There is no need to explore G. The subtree at G is pruned. Saves time.



The subtree below G is omitted

- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.
  - Case 1: happens on a Min node
  - Case 2: happens on a Max node

















- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.
  - Case 1: on a Min node; Case 2: on a Max node

slide 31



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do
  - β: the best (smallest) Min can do
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.

# **Alpha-beta pruning**

```
function Max-Value (s,\alpha,\beta)
inputs:
    s: current state in game, Max about to play
    \alpha: best score (highest) for Max along path to s
    \beta: best score (lowest) for Min along path to s
output: min(\beta, best-score (for Max) available from s)
    if (s is a terminal state)
    then return (terminal value of s)
    else for each s' in Succ(s)
     \alpha := \max(\alpha, Min-value(s', \alpha, \beta))
     if (\alpha \ge \beta) then return \beta /* alpha pruning */
    return a
function Min-Value(s,\alpha,\beta)
output: max(\alpha, best-score (for Min) available from s)
    if (s is a terminal state)
    then return (terminal value of s)
    else for each s' in Succs(s)
     \beta := \min(\beta, Max-value(s', \alpha, \beta))
     if (\alpha \ge \beta) then return \alpha /* beta pruning */
    return β
```

Starting from the root: Max-Value(root,  $-\infty$ ,  $+\infty$ )

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If at anytime  $\alpha$  exceeds  $\beta$ , the remaining children are pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.



- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_46_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_47_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_48_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_49_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_50_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_51_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_52_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_53_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_54_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_55_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_56_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_57_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_58_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_59_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_60_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_61_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_62_Figure_6.jpeg)

- Keep two bounds along the path
  - α: the best Max can do on the path
  - β: the best (smallest) Min can do on the path
- If a max node exceeds  $\beta$ , it is pruned.
- If a min node goes below  $\alpha$ , it is pruned.

![](_page_63_Figure_6.jpeg)

# How effective is alpha-beta pruning?

• Depends on the order of successors!

![](_page_64_Figure_2.jpeg)

- In the best case, the number of nodes to search is  $O(b^{m/2})$ , the square root of minimax's cost.
- This occurs when each player's best move is the leftmost child.
- In DeepBlue (IBM Chess), the average branching factor was about 6 with alpha-beta instead of 35-40 without.
- The worst case is no pruning at all.

# **Game-playing for large games**

- We've seen how to find game theoretic values. But it is too expensive for large games.
- What do real chess-playing programs do?
  - They can't possibly search the full game tree
  - They must respond in limited time
  - They can't pre-compute a solution

# **Game-playing for large games**

- The most popular solution: heuristic evaluation functions for games
  - 'Leaves' are intermediate nodes at a depth cutoff, not terminals
  - Heuristically estimate their values
  - Huge amount of knowledge engineering (R&N 6.4)
  - Example: Tic-Tac-Toe:

(number of 3-lengths open for me)-(number of 3-lengths open for you)

 Each move is a new depth-cutoff game-tree search (as opposed to search the complete game-tree once).

#### Prelude

 Heuristic estimation of the values of intermediate nodes can be done via Machine Learning

![](_page_67_Figure_2.jpeg)